

Network Working Group  
Request for Comments: 4905  
Category: Historic

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June 2007

## Encapsulation Methods for Transport of Layer 2 Frames over MPLS Networks

### Status of This Memo

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### Abstract

This document describes methods for encapsulating the Protocol Data Units (PDUs) of layer 2 protocols such as Frame Relay, Asynchronous Transfer Mode (ATM), or Ethernet for transport across an MPLS network. This document describes the so-called "draft-martini" protocol, which has since been superseded by the Pseudowire Emulation Edge to Edge Working Group specifications described in RFC 4447 and related documents.

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## 1. Introduction

In an MPLS network, it is possible to use control protocols such as those specified in [RFC4906] to set up "emulated virtual circuits" that carry the Protocol Data Units of layer 2 protocols across the network. A number of these emulated virtual circuits (VCs) may be carried in a single tunnel. This requires, of course, that the layer 2 PDUs be encapsulated. We can distinguish three layers of this encapsulation:

- the "tunnel header", which contains the information needed to transport the PDU across the MPLS network; this header belongs to the tunneling protocol, e.g., MPLS, Generic Routing Encapsulation (GRE), and Layer 2 Tunneling Protocol (L2TP).
- the "demultiplexer field", which is used to distinguish individual emulated virtual circuits within a single tunnel; this field must be understood by the tunneling protocol as well; it may be, e.g., an MPLS label or a GRE key field.
- the "emulated VC encapsulation", which contains the information about the enclosed layer 2 PDU that is necessary in order to properly emulate the corresponding layer 2 protocol.

This document specifies the emulated VC encapsulation for a number of layer 2 protocols. Although different layer 2 protocols require different information to be carried in this encapsulation, an attempt has been made to make the encapsulation as common as possible for all layer 2 protocols.

This document also specifies the way in which the demultiplexer field is added to the emulated VC encapsulation when an MPLS label is used as the demultiplexer field.

Quality of service (QoS)-related issues are not discussed in this document.

For the purpose of this document, R1 will be defined as the ingress router, and R2 as the egress router. A layer 2 PDU will be received at R1, encapsulated at R1, transported, decapsulated at R2, and transmitted out of R2.

## 2. Specification of Requirements

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

### 3. Special Note

This document describes the so called "draft-martini" protocol, which is used in many deployed implementations. This document and its contents have since been superseded by the Pseudowire Emulation Edge to Edge Working Group specifications: [RFC4447], [RFC4385], [RFC4448], [RFC4717], [RFC4618], [RFC4619], [RFC4553], [RFC4842], and related documents. This document serves as documentation of current implementations, and MUST NOT be used for new implementations. The PWE3 Label Distribution Protocol control protocol document [RFC4447], which is backward compatible with this document, MUST be used for all new implementations of this protocol.

### 4. General Encapsulation Method

In most cases, it is not necessary to transport the layer 2 encapsulation across the network; rather, the layer 2 header can be stripped at R1 and reproduced at R2. This is done using information carried in the control word (see below), as well as information that may already have been signaled from R1 to R2.

#### 4.1. The Control Word

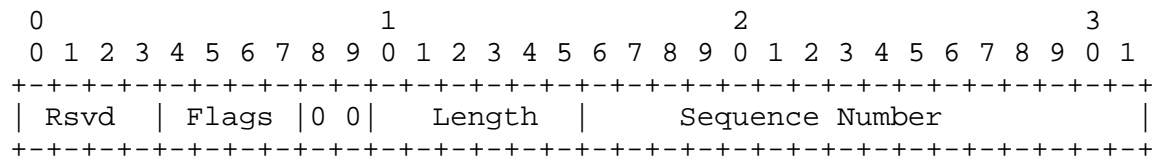
There are three requirements that may need to be satisfied when transporting layer 2 protocols over an MPLS backbone:

- i. Sequentiality may need to be preserved.
- ii. Small packets may need to be padded in order to be transmitted on a medium where the minimum transport unit is larger than the actual packet size.
- iii. Control bits carried in the header of the layer 2 frame may need to be transported.

The control word defined here addresses all three of these requirements. For some protocols, this word is REQUIRED, and for others OPTIONAL. For protocols where the control word is OPTIONAL, implementations MUST support sending no control word, and MAY support sending a control word.

In all cases, the egress router must be aware of whether the ingress router will send a control word over a specific virtual circuit. This may be achieved by configuration of the routers or by signaling, for example, as defined in [RFC4906].

The control word is defined as follows:



In the above diagram, the first 4 bits are reserved for future use. They MUST be set to 0 when transmitting, and MUST be ignored upon receipt.

The next 4 bits provide space for carrying protocol-specific flags. These are defined in the protocol-specific details below.

The next 2 bits MUST be set to 0 when transmitting.

The next 6 bits provide a length field, which is used as follows: If the packet's length (defined as the length of the layer 2 payload plus the length of the control word) is less than 64 bytes, the length field MUST be set to the packet's length. Otherwise, the length field MUST be set to 0. The value of the length field, if non-zero, can be used to remove any padding. When the packet reaches the service provider's egress router, it may be desirable to remove the padding before forwarding the packet.

The next 16 bits provide a sequence number that can be used to guarantee ordered packet delivery. The processing of the sequence number field is OPTIONAL.

The sequence number space is a 16-bit, unsigned circular space. The sequence number value 0 is used to indicate an unsequenced packet.

#### 4.1.1. Setting the Sequence Number

For a given emulated VC, and a pair of routers R1 and R2, if R1 supports packet sequencing, then the following procedures should be used:

- The initial packet transmitted on the emulated VC MUST use sequence number 1.
- Subsequent packets MUST increment the sequence number by 1 for each packet.
- When the transmit sequence number reaches the maximum 16 bit value (65535), the sequence number MUST wrap to 1.

If the transmitting router R1 does not support sequence number processing, then the sequence number field in the control word MUST be set to 0.

#### 4.1.2. Processing the Sequence Number

If a router R2 supports receive sequence number processing, then the following procedures should be used:

When an emulated VC is initially set up, the "expected sequence number" associated with it MUST be initialized to 1.

When a packet is received on that emulated VC, the sequence number should be processed as follows:

- If the sequence number on the packet is 0, then the packet passes the sequence number check.
- Else if the packet sequence number  $\geq$  the expected sequence number and the packet sequence number - the expected sequence number  $< 32768$ , then the packet is in order.
- Else if the packet sequence number  $<$  the expected sequence number and the expected sequence number - the packet sequence number  $\geq 32768$ , then the packet is in order.
- Otherwise, the packet is out of order.

If a packet passes the sequence number check or is in order, then it can be delivered immediately. If the packet is in order, then the expected sequence number should be set using the algorithm:

```
expected_sequence_number := packet_sequence_number + 1 mod 2**16
if (expected_sequence_number = 0) then expected_sequence_number := 1;
```

Packets that are received out of order MAY be dropped or reordered at the discretion of the receiver.

If a router R2 does not support receive sequence number processing, then the sequence number field MAY be ignored.

#### 4.2. MTU Requirements

The network MUST be configured with an MTU that is sufficient to transport the largest encapsulation frames. If MPLS is used as the tunneling protocol, for example, this is likely to be 12 or more bytes greater than the largest frame size. Other tunneling protocols may have longer headers and require larger MTUs. If the ingress

router determines that an encapsulated layer 2 PDU exceeds the MTU of the tunnel through which it must be sent, the PDU MUST be dropped. If an egress router receives an encapsulated layer 2 PDU whose payload length (i.e., the length of the PDU itself without any of the encapsulation headers) exceeds the MTU of the destination layer 2 interface, the PDU MUST be dropped.

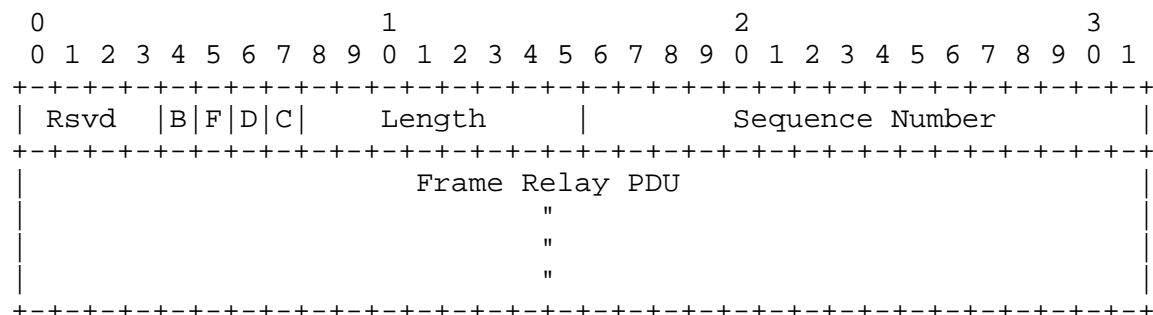
## 5. Protocol-Specific Details

### 5.1. Frame Relay

A Frame Relay PDU is transported without the Frame Relay header or the Frame Check Sequence (FCS). The control word is REQUIRED; however, its use is optional, although desirable. Use of the control word means that the ingress and egress Label Switching Routers (LSRs) follow the procedures below. If an ingress LSR chooses not to use the control word, it MUST set the flags in the control word to 0; if an egress LSR chooses to ignore the control word, it MUST set the Frame Relay control bits to 0.

The BECN (Backward Explicit Congestion Notification), FECN (Forward Explicit Congestion Notification), DE (Discard Eligibility), and C/R (Command/Response) bits are carried across the network in the control word. The edge routers that implement this document MAY, when either adding or removing the encapsulation described herein, change the BECN and/or FECN bits from 0 to 1 in order to reflect congestion in the network that is known to the edge routers, and the D/E bit from 0 to 1 to reflect marking from edge policing of the Frame Relay Committed Information Rate. The BECN, FECN, and D/E bits SHOULD NOT be changed from 1 to 0.

The following is an example of a Frame Relay packet:



\* B ( BECN ) Bit

The ingress router, R1, SHOULD copy the BECN field from the incoming Frame Relay header into this field. The egress router, R2, MUST generate a new BECN field based on the value of the B bit.

\* F ( FECN ) Bit

The ingress router, R1, SHOULD copy the FECN field from the incoming Frame Relay header into this field. The egress router, R2, MUST generate a new FECN field based on the value of the F bit.

\* D ( DE ) Bit

The ingress router, R1, SHOULD copy the DE field from the incoming Frame Relay header into this field. The egress router, R2, MUST generate a new DE field based on the value of the D bit.

If the tunneling protocol provides a field that can be set to specify a Quality of Service, the ingress router, R1, MAY consider the DE bit of the Frame Relay header when determining the value of that field. The egress router MAY then consider the value of this field when queuing the layer 2 PDU for egress. Note however that frames from the same VC MUST NOT be reordered.

\* C ( C/R ) Bit

The ingress router, R1, SHOULD copy the C/R bit from the received Frame Relay PDU to the C bit of the control word. The egress router, R2, MUST copy the C bit into the output frame.

## 5.2. ATM

Two encapsulations are supported for ATM transport: one for ATM Adaption Layer 5 (AAL5) and another for ATM cells.

The AAL5 Common Part Convergence Sublayer - Service Data Unit (CPCS-SDU) encapsulation consists of the REQUIRED control word and the AAL5 CPCS-SDU. The ATM cell encapsulation consists of an OPTIONAL control word, a 4-byte ATM cell header, and the ATM cell payload.

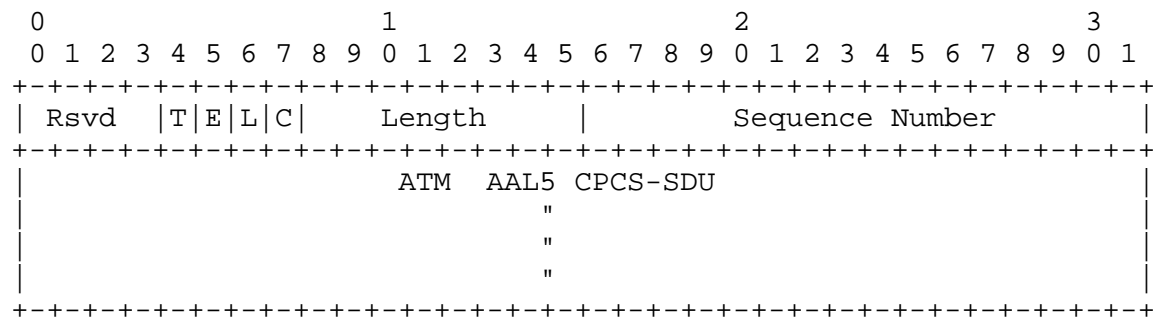


## 5.2.1. ATM AAL5 CPCS-SDU Mode

In ATM AAL5 mode, the ingress router is required to reassemble AAL5 CPCS-SDUs from the incoming VC and transport each CPCS-SDU as a single packet. No AAL5 trailer is transported. The control word is REQUIRED; its use, however, is optional, although desirable. Use of the control word means that the ingress and egress LSRs follow the procedures below. If an ingress LSR chooses not to use the control word, it MUST set the flags in the control word to 0; if an egress LSR chooses to ignore the control word, it MUST set the ATM control bits to 0.

The EFCI (Explicit Forward Congestion Indication) and CLP (Cell Loss Priority) bits are carried across the network in the control word. The edge routers that implement this document MAY, when either adding or removing the encapsulation described herein, change the EFCI bit from 0 to 1 in order to reflect congestion in the network that is known to the edge routers, and the CLP bit from 0 to 1 to reflect marking from edge policing of the ATM Sustained Cell Rate. The EFCI and CLP bits MUST NOT be changed from 1 to 0.

The AAL5 CPCS-SDU is prepended by the following header:



\* T (transport type) bit

Bit (T) of the control word indicates whether the packet contains an ATM cell or an AAL5 CPCS-SDU. If set, the packet contains an ATM cell, encapsulated according to the ATM cell mode section below; otherwise, it contains an AAL5 CPCS-SDU. The ability to transport an ATM cell in the AAL5 mode is intended to provide a means of enabling Operations and Management (OAM) functionality over the AAL5 VC.

\* E ( EFCI ) Bit

The ingress router, R1, SHOULD set this bit to 1 if the EFCI bit of the final cell of those that transported the AAL5 CPCS-SDU is set to 1, or if the EFCI bit of the single ATM cell to be transported in the packet is set to 1. Otherwise, this bit SHOULD be set to 0. The egress router, R2, SHOULD set the EFCI bit of all cells that transport the AAL5 CPCS-SDU to the value contained in this field.

\* L ( CLP ) Bit

The ingress router, R1, SHOULD set this bit to 1 if the CLP bit of any of the ATM cells that transported the AAL5 CPCS-SDU is set to 1, or if the CLP bit of the single ATM cell to be transported in the packet is set to 1. Otherwise, this bit SHOULD be set to 0. The egress router, R2, SHOULD set the CLP bit of all cells that transport the AAL5 CPCS-SDU to the value contained in this field.

\* C ( Command / Response Field ) Bit

When FRF.8.1 Frame Relay / ATM PVC Service Interworking [FRF.8.1] traffic is being transported, the CPCS-UU Least Significant Bit (LSB) of the AAL5 CPCS-SDU may contain the Frame Relay C/R bit. The ingress router, R1, SHOULD copy this bit to the C bit of the control word. The egress router, R2, SHOULD copy the C bit to the CPCS-UU Least Significant Bit (LSB) of the AAL5 CPCS PDU.

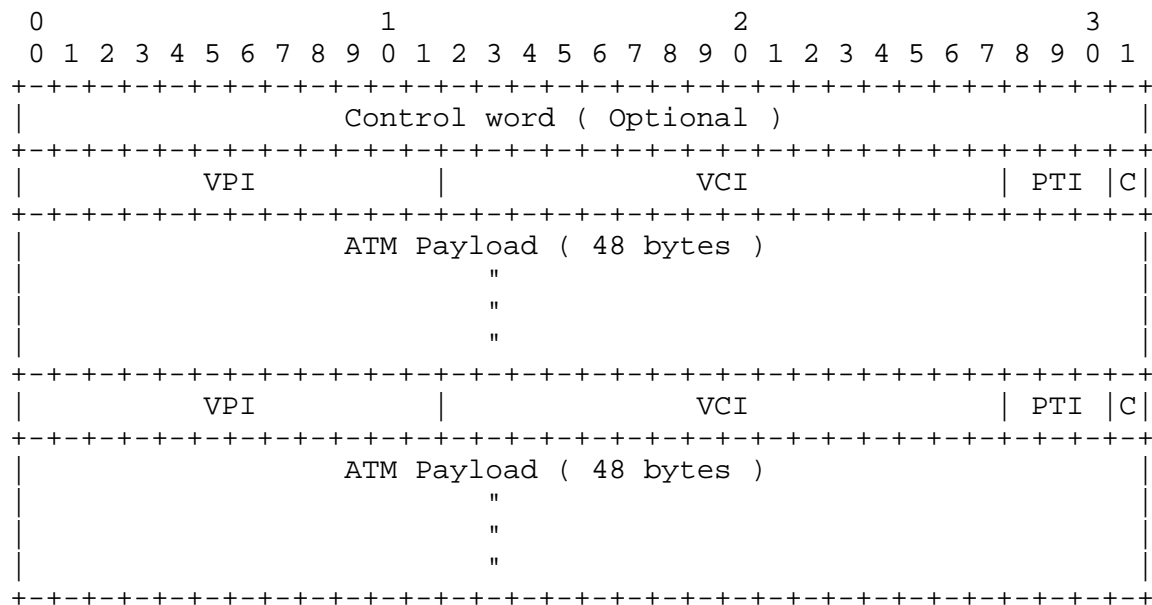
### 5.2.2. ATM Cell Mode

In this encapsulation mode, ATM cells are transported individually without a Segmentation and Reassembly (SAR) process. The ATM cell encapsulation consists of an OPTIONAL control word, and one or more ATM cells - each consisting of a 4-byte ATM cell header and the 48-byte ATM cell payload. This ATM cell header is defined in the FAST encapsulation [FAST] section 3.1.1, but without the trailer byte. The length of each frame, without the encapsulation headers, is a multiple of 52 bytes long. The maximum number of ATM cells that can be fitted in a frame, in this fashion, is limited only by the network MTU and by the ability of the egress router to process them. The ingress router MUST NOT send more cells than the egress router is willing to receive. The number of cells that the egress router is willing to receive may either be configured in the ingress router or may be signaled, for example, using the methods described in [RFC4906]. The number of cells encapsulated in a particular frame can be inferred by the frame length. The control word is OPTIONAL.

If the control word is used, then the flag bits in the control word are not used, and MUST be set to 0 when transmitting, and MUST be ignored upon receipt.

The EFCI and CLP bits are carried across the network in the ATM cell header. The edge routers that implement this document MAY, when either adding or removing the encapsulation described herein, change the EFCI bit from 0 to 1 in order to reflect congestion in the network that is known to the edge router, and the CLP bit from 0 to 1 to reflect marking from edge policing of the ATM Sustained Cell Rate. The EFCI and CLP bits SHOULD NOT be changed from 1 to 0.

This diagram illustrates an encapsulation of two ATM cells:



\* VPI (Virtual Path Identifier)

The ingress router **MUST** copy the VPI field from the incoming cell into this field. For particular emulated VCs, the egress router **MAY** generate a new VPI and ignore the VPI contained in this field.

\* VCI (Virtual Circuit Identifier)

The ingress router **MUST** copy the VCI field from the incoming ATM cell header into this field. For particular emulated VCs, the egress router **MAY** generate a new VCI.

\* PTI (Payload Type Identifier) & CLP ( C bit )

The PTI and CLP fields are the PTI and CLP fields of the incoming ATM cells. The cell headers of the cells within the packet are the ATM headers (without HEC) of the incoming cell.

### 5.2.3. OAM Cell Support

OAM cells MAY be transported on the VC LSP. An egress router that does not support transport of OAM cells MUST discard frames that contain an ATM cell with the high-order bit of the PTI field set to 1. A router that supports transport of OAM cells MUST follow the procedures outlined in [FAST] section 8 for mode 0 only, in addition to the applicable procedures specified in [RFC4906].

### 5.2.4. CLP bit to Quality of Service Mapping

The ingress router MAY consider the CLP bit when determining the value to be placed in the Quality of Service fields (e.g., the EXP fields of the MPLS label stack) of the encapsulating protocol. This gives the network visibility of the CLP bit. Note however that cells from the same VC MUST NOT be reordered.

### 5.3. Ethernet VLAN

For an Ethernet 802.1q VLAN, the entire Ethernet frame without the preamble or FCS is transported as a single packet. The control word is OPTIONAL. If the control word is used, then the flag bits in the control word are not used, and MUST be set to 0 when transmitting, and MUST be ignored upon receipt. The 4-byte VLAN tag is transported as is, and MAY be overwritten by the egress router.

The ingress router MAY consider the user priority field [IEEE802.3ac] of the VLAN tag header when determining the value to be placed in the Quality of Service field of the encapsulating protocol (e.g., the EXP fields of the MPLS label stack). In a similar way, the egress router MAY consider the Quality of Service field of the encapsulating protocol when queuing the packet for egress. Ethernet packets containing hardware-level Cyclic Redundancy Check (CRC) errors, framing errors, or runt packets MUST be discarded on input.

### 5.4. Ethernet

For simple Ethernet port to port transport, the entire Ethernet frame without the preamble or FCS is transported as a single packet. The control word is OPTIONAL. If the control word is used, then the flag bits in the control word are not used, and MUST be set to 0 when transmitting, and MUST be ignored upon receipt. As in the Ethernet

VLAN case, Ethernet packets with hardware-level CRC errors, framing errors, and runt packets MUST be discarded on input.

### 5.5. High-Level Data Link Control (HDLC)

HDLC mode provides port to port transport of HDLC-encapsulated traffic. The HDLC PDU is transported in its entirety, including the HDLC address, control, and protocol fields, but excluding HDLC flags and the FCS. Bit/byte stuffing is undone. The control word is OPTIONAL. If the control word is used, then the flag bits in the control word are not used, and MUST be set to 0 when transmitting, and MUST be ignored upon receipt.

The HDLC mode is suitable for port to port transport of Frame Relay User-Network Interface (UNI) or Network-Network Interface (NNI) traffic. It must be noted, however, that this mode is transparent to the FECN, BECN, and DE bits.

### 5.6. PPP

PPP mode provides point to point transport of PPP-encapsulated traffic, as specified in [RFC1661]. The PPP PDU is transported in its entirety, including the protocol field (whether compressed using PFC or not), but excluding any media-specific framing information, such as HDLC address and control fields or FCS. Since media-specific framing is not carried, the following options will not operate correctly if the PPP peers attempt to negotiate them:

- Frame Check Sequence (FCS) Alternatives
- Address-and-Control-Field-Compression (ACFC)
- Asynchronous-Control-Character-Map (ACCM)

Note also that VC LSP Interface MTU negotiation as specified in [RFC4906] is not affected by PPP Maximum Receive Unit (MRU) advertisement. Thus, if a PPP peer sends a PDU with a length in excess of that negotiated for the VC LSP, that PDU will be discarded by the ingress router.

The control word is OPTIONAL. If the control word is used, then the flag bits in the control word are not used, and MUST be set to 0 when transmitting, and MUST be ignored upon receipt.

## 6. Using an MPLS Label as the Demultiplexer Field

To use an MPLS label as the demultiplexer field, a 32-bit label stack entry [RFC3032] is simply prepended to the emulated VC encapsulation, and hence will appear as the bottom label of an MPLS label stack. This label may be called the "VC label". The particular emulated VC

identified by a particular label value must be agreed by the ingress and egress LSRs, either by signaling (e.g., via the methods of [RFC4906]) or by configuration. Other fields of the label stack entry are set as follows.

#### 6.1. MPLS Shim EXP Bit Values

If it is desired to carry Quality of Service information, the Quality of Service information SHOULD be represented in the EXP field of the VC label. If more than one MPLS label is imposed by the ingress LSR, the EXP field of any labels higher in the stack SHOULD also carry the same value.

#### 6.2. MPLS Shim S Bit Value

The ingress LSR, R1, MUST set the S bit of the VC label to a value of 1 to denote that the VC label is at the bottom of the stack.

#### 6.3. MPLS Shim TTL Values

The ingress LSR, R1, SHOULD set the TTL field of the VC label to a value of 2.

### 7. Security Considerations

This document specifies only encapsulations, and not the protocols, used to carry the encapsulated packets across the network. Each such protocol may have its own set of security issues, but those issues are not affected by the encapsulations specified herein. More detailed security considerations are also described in Section 8 of [RFC4447].

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## Acknowledgement

Funding for the RFC Editor function is currently provided by the Internet Society.

