

RIP Version 2 Protocol Applicability Statement

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

Abstract

As required by Routing Protocol Criteria (RFC 1264), this report defines the applicability of the RIP-2 protocol within the Internet. This report is a prerequisite to advancing RIP-2 on the standards track.

1. Protocol Documents

The RIP-2 protocol analysis is documented in RFC 1721 [1].

The RIP-2 protocol description is defined in RFC 1723 [2]. This memo obsoletes RFC 1388, which specifies an update to the "Routing Information Protocol" RFC 1058 (STD 34).

The RIP-2 MIB description is defined in RFC 1724 [3]. This memo will obsolete RFC 1389.

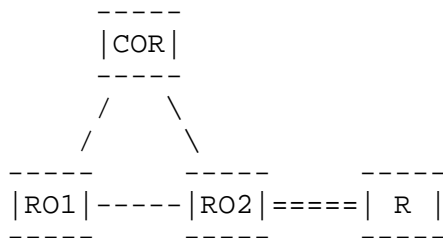
2. Introduction

This report describes how RIP-2 may be useful within the Internet. In essence, the environments in which RIP-2 is the IGP of choice is a superset of the environments in which RIP-1, as defined in RFC 1058 [1], has traditionally been used. It is important to remember that RIP-2 is an extension to RIP-1; RIP-2 is not a new protocol. Thus, the operational aspects of distance-vector routing protocols, and RIP-1 in particular, within an autonomous system are well understood.

It should be noted that RIP-2 is not intended to be a substitute for OSPF in large autonomous systems; the restrictions on AS diameter and complexity which applied to RIP-1 also apply to RIP-2. Rather, RIP-2 allows the smaller, simpler, distance-vector protocol to be used in environments which require authentication or the use of variable

as a next hop, even if it is a better next hop than XR1. Of course, XR1 knows this and can indicate, via the Next Hop field, that XR2 is the better next hop for some routes.

Another use for Next Hop has also been found. Consider the following example topology:



The three links between the Central Office Router (COR) and the Remote Office routers (R01 and R02) are all Dial-On-Demand (DOD) links. The link between R02 and R is a fixed link. Once all of the routers have been initialized, the only routes they know about are the configured static routes for the DOD links. Assume that connections between COR and R01, and COR and R02 are established and RIP information is passing between the routers. R01 will ignore COR's route to R02 because it already has a better one; however, it will learn to reach R via COR.

If we assume that R01 and R02 are only capable of establishing one link at a time, then R01 will not be able to reach R02; however, R01 will be able to reach R. Worse still, if we assume that traffic stops and the DOD links drop due to inactivity, an attempt by R01 to reach R will trigger the dialing of two links (through COR). Of course, once R01 establishes a link to R02, the problem corrects itself because the new route to R is one hop shorter.

To correct this problem, the routers may use the Next Hop field to indicate their next hop. Consider the following route advertisements during the period described above (before the R01/R02 link has ever been established):

Sender	Recvr	Route	NextHop	Metric
=====				
R02	COR	R	0	1

COR	R01	R02	0	1
		R	R02	2

When R01 receives the two routes from COR, it will ignore the route for R02, as mentioned above. However, since R is not in R01's routing table, it will add it using a next hop of R02 (because R02 is directly connected, after a fashion). Note that COR does count itself in R's metric; this is less than accurate, but entirely safe and correctable (when the R01/R02 link comes up). Suppose, now, that the R01/R02 link did not exist. R01 would ignore the specification of R02 as the next hop to R and use COR, as it would if no Next Hop had been specified.

Note that this is not a recursive algorithm; it only works to eliminate a single extra hop from the path. There are methods by which this mechanism might be extended to include larger optimizations, but the potential to create routing loops has not been sufficiently analyzed to specify them here.

3.3 Authentication

The need for authentication in a routing protocol is obvious. It is not usually important to conceal the information in the routing messages, but it is essential to prevent the insertion of bogus routing information into the routers. So, while the authentication mechanism specified in RIP-2 is less than ideal, it does prevent anyone who cannot directly access the network (i.e., someone who cannot sniff the routing packets to determine the password) from inserting bogus routing information.

However, the specification does allow for additional types of authentication to be incorporated into the protocol. Unfortunately, because of the original format of RIP packets, the amount of space available for providing authentication information is only 16 octets.

3.4 Multicasting

The RIP-2 protocol provides for the IP multicasting of periodic advertisements. This feature was added to decrease the load on systems which do not support RIP-2. It also provides a mechanism whereby RIP-1 routers will never receive RIP-2 routes. This is a feature when correct use of an advertised route depends on knowing the precise subnet mask, which would be ignored by a RIP-1 router.

4. Conclusion

Because the basic protocol is unchanged, RIP-2 is as correct a routing protocol as RIP-1. The enhancements make RIP-2 useful in environments which RIP-1 could not handle, but which do not necessitate the use of OSPF by virtue of requirements which RIP-2 does not satisfy.

5. References

- [1] Malkin, G., "RIP Version 2 Protocol Analysis", RFC 1721, Xylogics, Inc., November 1994.
- [2] Malkin, G., "RIP Version 2 - Carrying Additional Information", RFC 1723, Xylogics, Inc., November 1994.
- [3] Malkin, G., and F. Baker, "RIP Version 2 MIB Extension", RFC 1724, Xylogics, Inc., Cisco Systems, November 1994.

6. Security Considerations

Security issues are not discussed in this memo.

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